data egret



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Plan

- What is a transaction and why do we need them
- How transactions are implemented
- Transactions in PostgreSQL



Why do we need transactions?

How much?

```
Sum: EUR 1000

to account B -100€
to account C -200€

Sum: EUR ?
```

How?

```
send_money(src_acc, dst_acc, amount):
    balance := src_acc.balance();
    if( balance - amount > 0 ):
        dst_acc.balance += amount;
        src_acc.balance = balance - amount;
        return 0;
    else:
        return 1;
```



Why do wee need transactions?

Action 1

```
send_money(src_acc, dst_acc, amount):
    balance := src_acc.balance();
    if( balance - amount > 0 ):
        dst_acc.balance += amount;

        src_acc.balance = balance - amount;
        return 0;
    else:
        return 1;
```

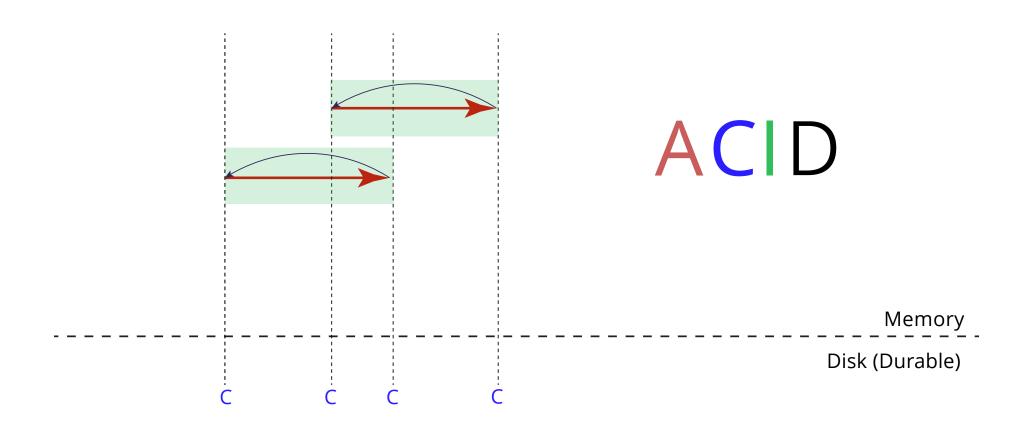
Action 2

```
send_money(src_acc, dst_acc, amount):
    balance := src_acc.balance();
    if( balance - amount > 0 ):
        dst_acc.balance += amount;
        src_acc.balance = balance - amount;
        return 0;
    else:
        return 1;
```

What do we already learned?

- Transfer actions consist of smaller simple actions
- Results might be affected by order
- If there is only one action, we do not have any problems
- There is no problem to read data
- Problems come when we write data
- More concurrency lead to more problems

What can we do?





Action with ACID properties

- Atomicity happens completely or fails completely
- Consistency brings data from one cosistent state to another
- Isolation "thinks" it happens alone
- Durability what is saved to disk is safe

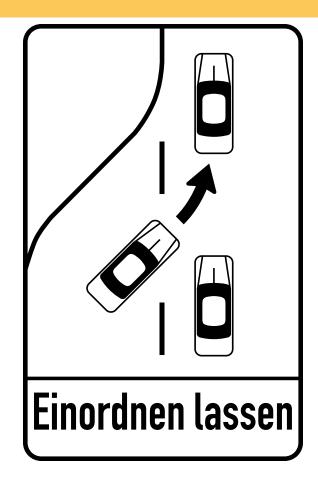
Action with ACID properties

- Atomicity happens completely or fails completely
- Consistency brings data from one cosistent state to another
- Isolation "thinks" it happens alone
- **Durability** what is saved to disk is safe

We call such an action a Transaction



How transactions are implemented?



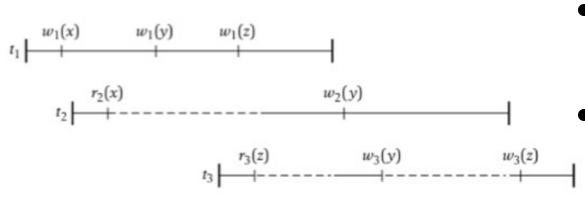
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Problem: Serializability

- In theory everything is simple
- ... as Reißverschlussverfahren
- But in practice it is either slow or cause failures
- ...or we need an algorythm



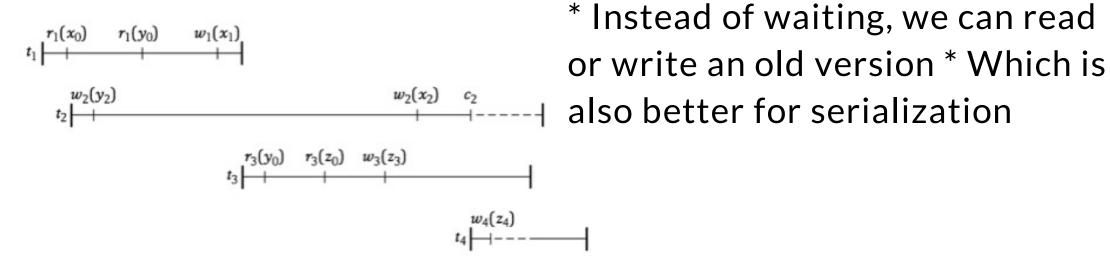
2 Phase Locking



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- 1. Phase: We take all locks this scope of transactions needs
- 2. Phase We release all those
 locks and do not take new locks
- It works, but:
 - it is slow
 - Deadlocks

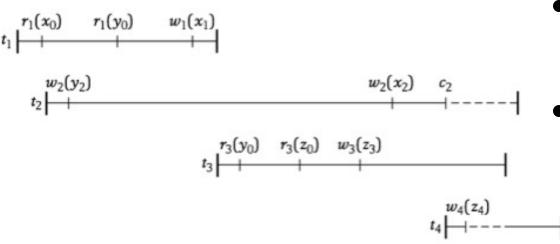
Multi Version 2 Phase Locking



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Multi Version 2 Phase Locking

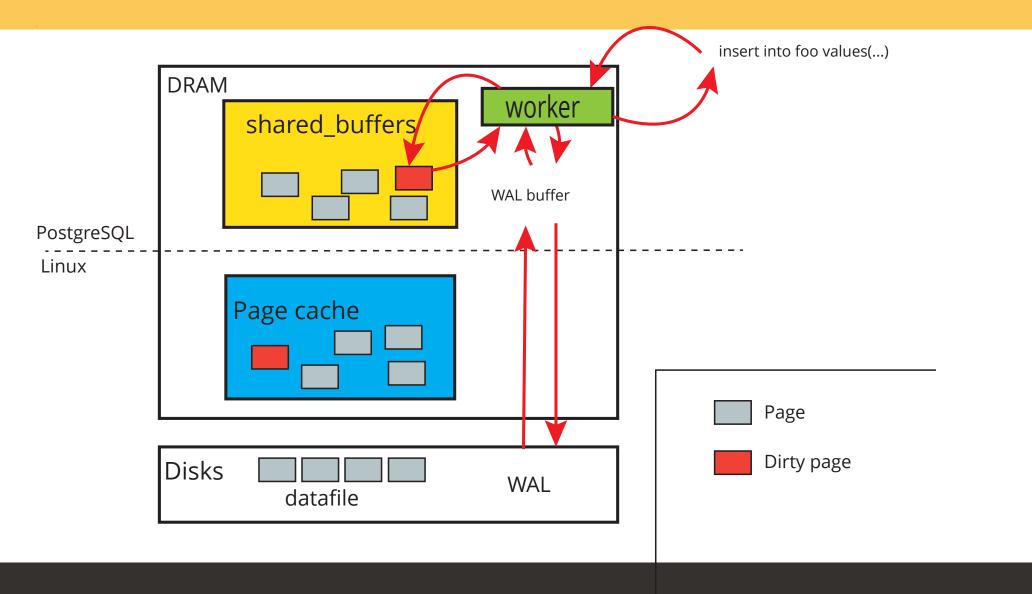


- Instead of waiting, we can read or write an old version
- Which is also better for serialization

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Postgres: inserting data



Postgres: inserting data

Internal

```
StartTransactionCommand;
    StartTransaction;
ProcessUtility;
    BeginTransactionBlock;
CommitTransactionCommand;
StartTransactionCommand;
ProcessQuery;
CommitTransactionCommand;
    CommandCounterIncrement;
StartTransactionCommand;
ProcessUtility;
    EndTransactionBlock;
CommitTransactionCommand;
    CommitTransaction;
```

SQL

```
BEGIN;
INSERT...
COMMIT
```

in PostgreSQL everything is transactional

- a writing transaction get a XID
- each tuple has xmin (XID of the "yangest" transaction which has updated this tuple
- and xmax (XID of the oldest transaction which can see this tuple
- each backend has its xmin and syncronizes it through
 MyProc->xmin it is a way how Snapshot works



PostgreSQL is a MVCC Database

- that dooesn't mean, that there is no 2PL
- it is possible to go through many old versions, but through all of them
- At some point we were considering such an architeckture as a fail, but now it is rather future

Questions?

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